

Percolation theory on multirooted directed graphs

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The multiroot connectedness P_{uv} of a directed graph G between the vertex u and the collection of vertices $\mathbf{v} = \{v_1, \dots, v_n\}$ is the probability that there are (directed) paths from u to each of the vertices v_i , $i = 1, \dots, n$, when each edge and vertex has a given probability of being independently deleted. The properties of the coefficients in the expansion

$$P_{uv}(G) = \sum_{A' \subseteq A} \vec{d}_{uv}(G') \prod_{a \in A'} p_a \prod_{w \in V} p_w,$$

where A and V are the arc and vertex sets of G respectively, p_a (p_w) is the probability that an arc (vertex w) is not deleted, and G' is the arc set A' together with its incident vertices V' , are considered. The values of $\vec{d}_{uv}(G)$ are characterized as follows: $\vec{d}_{uv}(G)$ is shown to be nonzero if and only if G is coverable by paths and has no directed circuit. For this case $\vec{d}_{uv}(G) = (-1)^{t_{uv} + n} = (-1)^{\nu(G)}$, where t_{uv} is the maximal number of independent directed paths between u and the set \mathbf{v} , and t_{uv} is shown to be equal to $\nu(G) + n$, where $\nu(G) = |E| - |V| + 1$ is the cyclomatic number of G .

INTRODUCTION

This paper generalizes the results in Ref. 1 on two rooted directed graphs to directed graphs with many roots. The weights associated with such graphs and their properties proved here are required in a forthcoming paper² on percolation theory in a gas. The result $\vec{d}_{uv}(G) = (-1)^{\nu(G)}$ is also used in developing a relation between graphical expansions and renormalization for the percolation problem.³

Many of the proofs of the results given here follow those given in Ref. 1 and so they are omitted. Only the differences will be emphasized between the many-root and two-rooted situations.

A completely different proof of Theorem 5 on the directed-undirected d weight relation is given⁴ in contrast to the inductive proof in Ref. 1.

Consider a graph G with input vertex u and output vertices $\mathbf{v} = \{v_1, \dots, v_n\}$ and suppose that a subset of the edges and vertices is deleted. G is sometimes denoted by G_{uv} if we wish to emphasize the root points of G . Let $S_{u,v}$ be the set of all self-avoiding paths (following the arrows if G is directed) from u to v_i on G . For $s \in S_{u,v}$, define the indicator random variable

$$\gamma(s) = \begin{cases} 1, & \text{if } s \text{ is open,} \\ 0, & \text{if } s \text{ is closed,} \end{cases}$$

in a given state of multilation of G . Now

$$\gamma_i \equiv \sum_{\phi \subset S_i \subseteq S_{u,v}} (-1)^{|S_i|+1} \prod_{s \in S_i} \gamma(s) \quad (2.1)$$

is one or zero according as there is at least one open path or no open path from u to v_i . The pair-connectedness is therefore given between u and v_i by

$$P_{u,v_i} = \langle \gamma_i \rangle = \sum_{\phi \subset S_i \subseteq S_{u,v}} (-1)^{|S_i|+1} \langle \prod_{s \in S_i} \gamma(s) \rangle.$$

Now let $g(S_i)$, $V(S_i)$, $E(S_i)$, $A(S_i)$ be the graph, the vertex set, the edge set, and the arc set obtained by taking the union of the paths in S_i . Thus if p_a is the probability that element α is not deleted, then

$$\langle \prod_{s \in S_i} \gamma(s) \rangle = \prod_{v \in V(S_i)} p_v \cdot \prod_{e \in E(S_i)} p_e,$$

since this is the probability that all the paths in S_i are open. The result may be expressed as a sum over subgraphs of G by grouping together all S_i for which $E(S_i) = E'$, and $g(S_i) = G'$; thus

$$P_{u,v_i} = \sum_{E' \subseteq E} d_{u,v_i}(G') \prod_{v \in V'} p_v \prod_{e \in E'} p_e,$$

where the " d weight" is given by

$$d_{u,v_i}(G') = \sum_{\substack{\phi \subset S_i \subseteq S_{u,v} \\ g(S_i) = G'}} (-1)^{|S_i|+1}.$$

For the directed case read A for E and \vec{d} for d .

The extension to many outputs is now straightforward.

$P_{u,v}$ = probability that there is a path from u to all of the output vertices

$$= \langle \prod_{i=1}^n \gamma_i \rangle = \sum_{E' \subseteq E} d_{u,v}(g') \prod_{v \in V'} p_v \prod_{e \in E'} p_e \quad (2.2)$$

where

$$d_{u,v}(G') = \prod_{i=1}^n \sum_{\phi \subset S_i \subseteq S_{u,v}} (-1)^{|S_i|+1} \quad (2.3)$$

$$g\left(\bigcup_{i=1}^n S_i\right) = G'$$

If there does not exist $S = \bigcup_{i=1}^n S_i$ such that $g(S) = G'$, then $\vec{d}_{u,v}(G') = 0$ from (2.3).

To establish a contraction-deletion rule for $\vec{d}_{u,v}$ weights let $p_a = 1$, $p_w = 1$ in (2.2). Then

$$\vec{\gamma}_{u,v}(G) = \sum_{A' \subseteq A} \vec{d}_{u,v}(G'), \quad (2.4)$$

where the connectedness indicator $\vec{\gamma}_{u,v}(G)$ is 1 if there are paths from u to each v_i of $G_{u,v}$ and zero otherwise.

Since the set of all subsets of the arc set A form a lattice, (2.4) may be inverted² to give

$$\vec{d}_{uv}(G) = \sum_{A' \subseteq A} (-1)^{|A \setminus A'|} \vec{\gamma}_{uv}(G'). \quad (2.5)$$

Now consider a particular arc $a \in A$ and divide the sum according as $a \in A'$ and $a \notin A'$. We get

$$\vec{d}_{uv} \sum_{\substack{A' \subseteq A \\ a \in A'}} (-1)^{|A \setminus A'|} \vec{\gamma}_{uv}(G') + \sum_{\substack{A' \subseteq A \\ a \notin A'}} (-1)^{|A \setminus A'|} \vec{\gamma}_{uv}(G'). \quad (2.6)$$

Let G^γ be the graph G contracted along the arc a ; that is, G^γ is the graph $G \setminus \{a\}$ with the adjacent vertices of a identified. Also we define G^δ to be the graph G with arc a deleted. If G' is a subgraph of G , we define G'^γ to be the graph with the arc and vertex set of G' with arc a contracted, and G'^δ to be G' less the arc a .

Thus if we ensure the arc a is oriented out of u it is clear that $\vec{\gamma}_{uv}(G') = \vec{\gamma}_{uv}(G'^\gamma)$ when $a \in A'$. This true even if the arc a has a root v_i as an adjacent vertex if we assume that when $u = v_i$, then u and v_i are connected. Furthermore in the second summation G' does not contain the arc a and so $\vec{\gamma}_{uv}(G') = \vec{\gamma}_{uv}(G'^\delta)$ since $G' = G'^\delta$.

It follows that

$$\vec{d}_{uv}(G) = \vec{d}_{uv}(G^\gamma) - \vec{d}_{uv}(G^\delta) \quad (2.7)$$

subject to a being oriented out of u .

Remark: When $u = v_i$ on contraction, $\vec{d}_{uv}(G^\gamma) = \vec{d}_{uv'}(G^\gamma)$, where $v' = v \setminus \{v_i\}$ if we define $\vec{\gamma}_{uu}(G^\gamma) = 1$.

2. STATEMENT OF RESULTS

As stated above the d weight of a graph which is not coverable is zero. The following theorems apply to the d weight of a coverable multirooted graph G .

Theorem 1: The directed d weight is ± 1 or 0.

Theorem 2: The directed d weight is zero if and only if G has a circuit.

Definition: A collection $\mathcal{C} = \{\pi_i | i = 1, \dots, n\}$ of (directed) paths on a coverable directed graph G_{uv} is said to be independent if the matrix $M = [m_{ij}]$ has maximal row rank where $\pi_i = \sum_{j=1}^l m_{ij} a_j$, $\pi_i \in \mathcal{C}$, and $a_{ij} = 1, \dots, l$ is the collection of arcs in G_{uv} , and $m_{ij} \in \mathbb{Q}$. Such a collection \mathcal{C} is said to be maximal if every path not in \mathcal{C} is dependent on the elements of \mathcal{C} . Obviously the number of paths in such a class is an invariant of G_{uv} . We will call such a set a maximal independent set (MIS) for G_{uv} .

Theorem 3: If the directed graph has no circuit, then

$$\vec{d}_{uv}(G) = (-1)^{t_{uv} + n}$$

where t_{uv} is the order of a MIS for G_{uv}

Theorem 4: For a directed graph G with no circuit $t_{uv} = \nu(G) + n$, where $\nu(G) = |E| - |V| + 1$ is the cyclomatic number for G .

Note $\vec{d}_{uv}(G) = (-1)^{\nu(G)}$ follows from Theorems 3 and 4.

Theorem 5: The undirected d weight of G is equal to the

sum of the directed \vec{d} weights of G over all possible orientations of G .

3. Remarks and proofs of results

The basic graph which is the terminal stage of the reduction process using the contraction–deletion rule which was described in Sec. 1 is the multiroot parallel graph: Every arc is adjacent to the input root point u and one of the output root points of v .

Proof of Theorem 1: First of all it is assumed that v , the set of output root points, contains sinks and at each stage in the reduction process the arc $[u, w]$ is chosen so that w is not a sink. The theorem then follows since the reduction process ensures that $\vec{d}_{uv}(G) = 0$ or $|\vec{d}_{uv}(G)| = |\vec{d}_{uv}(G')|$, where G' is a parallel graph with input u and output $v_k = \{v_i, \dots, v_k\}$, the set of sinks in v . Therefore, $|\vec{d}_{uv}(G)| = \pm 1$ or 0.

For the purposes of calculation the assumption that $v_k \neq \emptyset$ is no real restriction because $\vec{d}_{uv}(G) = \vec{d}_{uv}(G_e)$ where G_e is the union of G with a single arc b attached to a root point v_i where v_i is replaced by v_e as the root point, and b is oriented from v_i to v_e .

It can be further seen that if G does not have any root points which are sinks then there is a directed circuit in G and so the first part of Theorem 2 applies which will result in $\vec{d}_{uv}(G) = 0$.

Proof of Theorem 2: This proof follows the corresponding Theorem 3 in Ref. 1. The only essential difference is that if $[u, w]$, the arc upon which the contraction–deletion rule is applied, is such that w is a root, then

$$\{\pi_i\}_1 \cup \{\pi_i\}_1 \cup \{\pi_i \circ \pi_j\}_1^t$$

is a covering for G^δ .

Proof of Theorem 3: As in the proof of the first part of Theorem 2 (\Rightarrow), we can apply the contraction–deletion to an arc a of G with vertices u and w and contain two cases:

- (1) $\vec{d}_{uv}(G^\gamma) = \vec{d}_{uv}(G)$, G^γ has no circuit,
- (2) $\vec{d}_{uv}(G^\delta) = -\vec{d}_{uv}(G)$, G^δ has no circuit,

where G^γ and G^δ are both coverable. When w is not a root point it is easily shown, as in Theorem 4,¹ that number of independent paths in G and G^γ are equal and differs from the number independent paths in G^δ by one.

If w is a root point say v_1 , then there is a natural one–one correspondence between paths on G^γ and the set of paths on G not including the root path π consisting of the (contracted) arc a between u and v_1 .

A maximal independent set of contracted paths \mathcal{C}^γ on G^γ gives rise to an independent set \mathcal{C}' on G using the 1–1 correspondence.

We claim that $\mathcal{C} = \mathcal{C}' \cup \{\pi\}$ is a maximal independent set for paths on G . The set \mathcal{C} is independent because if $\pi = \sum_{i=1}^n \alpha_i \pi_i$, $\pi_i \in \mathcal{C}'$, $\alpha_i \in \mathbb{Q}$, then on the contracted graph G^γ , $\pi^\gamma = \sum_{i=1}^n \alpha_i \pi_i^\gamma$; but $\pi^\gamma = \mathbf{0}$, the π_i^γ are independent, and not all the α_i 's are zero since $\pi \neq \mathbf{0}$, so we have a contradiction.

To show that \mathcal{C} is maximal on G suppose there exists a path $\pi_0 (\neq \pi)$ such that $\mathcal{C} \cup \{\pi_0\}$ is independent. However the contracted path π_0^γ must be a linear sum of the MIS for G^γ , that is

$$\pi_0^\gamma = \sum_{i=1}^n \alpha_i \pi_i^\gamma, \pi_i^\gamma \in C^\gamma, \alpha_i \in \mathbb{Q}, \alpha_i \neq 0, \text{ for some } i.$$

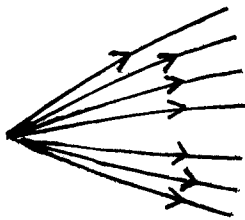
For the cases $a \in \pi_0, a \notin \pi_0$ the coefficients α_i at the vertex w sum to 1 and 0, respectively, therefore

$$\pi_0 = \sum_{i=1}^n \alpha_i \pi_i, \text{ and so we have a contradiction.}$$

Thus when w is a root point the number of independent paths in G and G^γ differ by one.

The other cases follow as in Theorem 4.¹

Repeated application of the contraction-deletion process results in a parallel graph H of the following type:



Note there will in general be only k of the original n roots, say $v' = \{v_1, \dots, v_k\}$, the set of sinks which is nonempty by the no-circuit property. Moreover $d_{uv}(H) = +1 = (-1)^{t_{uv} + k}$, where $t_{uv} (=k)$ is the number of independent paths in H .

Furthermore in the reduction process from G to H the number of independent paths has changed by one: (a) for every reduction by deletion, (b) for every reduction by contraction of a root point.

Therefore, the total change in the number of independent paths $t_{uv} - t_{uv_k} = (n - k) + \#$ of deletions

and hence the number of sign changes in the \vec{d} weight for the sequence of graphs from G to H is

$$t_{uv} - t_{uv_k} - (n - k),$$

since $n - k$ is the number of roots contracted.

Therefore,

$$\begin{aligned} \vec{d}_{uv}(G) &= (-1)^{t_{uv_k} + k + t_{uv} - t_{uv_k} - (n - k)} \\ &= (-1)^{t_{uv} + n}, \end{aligned}$$

where t_{uv} is the number of independent paths in G .

Proof of Theorem 4: Let $\Pi = \{\pi_i | i = 1, \dots, l\}$ be a MIS for G . The generalization required here is that G_0 be a subgraph of G consisting of a subset $\{\pi_i | i = 1, \dots, n\}$ of Π covering all the root points. We can always take G_0 to be a tree.

The order of a MIS for G_0 is $n = |v|$. Moreover it is easy to check that $\vec{d}_{uv}(G) = +1$. Let ρ_1, \dots, ρ_k be paths in Π which are not in G_0 .

Define a sequence of subgraphs $G_i = G_0 \cup \{\cup_{j=1}^i \rho_j\}$. For the sequence G_i define $\mu(G_i) = |E_i| - |V_i| + 1 + n$. We note $\mu(G_0) = n$ which is the number of paths in a MIS for G_0 .

The argument now closely follows Theorem 5¹ and we obtain $\mu(G) = \mu(G_k) = t_{uv}$. Now $\mu(G) = \nu(G) + n$, where $\nu(G)$ is the cyclomatic number.⁵ Therefore,

$$\vec{d}_{uv}(G) = (-1)^{t_{uv} + n} = (-1)^{\nu(G) + n + n} = (-1)^{\nu(G)}.$$

Proof of Theorem 5: Let G be an undirected graph and $\mathcal{D}(G)$ be the set of graphs obtained by directing G in all possible ways. Now consider the paths on G as directed from u to v . Any subset of \mathcal{S}_i which covers G either covers some $g \in \mathcal{D}(G)$ or $g(S_i)$ has at least one loop of length 2. Let $\mathcal{L}(G)$ be the set of graphs obtained from g by directing it in all possible ways and replacing at least one edge by a loop of length 2. Thus

$$\begin{aligned} d_{uv}(G) &= \sum_{g \in \mathcal{D}(G)} \sum_{\substack{\phi \subset S_i \subseteq S_{uv} \\ g(S_i) = \vec{g}}} (-1)^{|S_i| + 1} \\ &\times \sum_{g \in \mathcal{L}(G)} \sum_{\substack{\phi \subset S_i \subseteq S_{uv} \\ g(S_i) = \vec{g}}} (-1)^{|S_i| + 1}. \end{aligned} \quad (4.1)$$

But for $\vec{g} \in \mathcal{L}(G)$,

$$\sum_{\substack{\phi \subset S_i \subseteq S_{uv} \\ g(S_i) = \vec{g}}} (-1)^{|S_i| + 1} = 0 \quad (4.2)$$

since it is the \vec{d} weight of a directed graph with a loop and the result follows by substituting (4.1) with condition (4.2) in (2.3).

¹D. K. Arrowsmith and J. W. Essam, "Percolation theory on directed graphs," *J. Math. Phys.* **18**, 235-8 (1977).

²J. W. Essam and A. Coniglio, "Percolation Theory in a Gas," *J. Phys. A: Math. Gen.* **10**, 1917-26 (1977).

³J. W. Essam and C. M. Place, "Low Density Expansion of the Pair Connectedness for Percolation Models," *Ann. Israel. Phys. Soc.* **2**, 882 (1978).

⁴Private communication from J. W. Essam.

⁵C. Berge, *The Theory of Graphs* (Methuen, London, 1962).